The Internet Architecture [Clark88a]

CSci551: Computer Networks SP2006 Thursday Section John Heidemann

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Key ideas

- · motivation behind Internet architecture
 - why things are they way they are
 - hints and some alternative designs (and why they weren't taken)
- not complete

 - main goal of internet
 building on existing networks
 - secondary goals
 - · reliability / availability
 - approaches
 - starting with simple building blocks
 - datagram

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What is the Internet?

- *inter*net: composition of many small neteworks
- a common network protocol IP
- multiple protocols: TCP, UDP, etc.
- multiple services: HTTP (web), SMTP (mail)
- dynamic routing protocols
 - routes can change automatic (different links)
 - routes change do to network conditions
 - change due to SOME network conditions - link failures
 - evalute and change to congestion infrequently: traffic engineering
 - · not so much due to congestion

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The Internet Architecture

- primary goals:
 - connecting heterogeneous networks
 - packet switching
- secondary goals:
 - robusteness
 - (rest are in list in paper)

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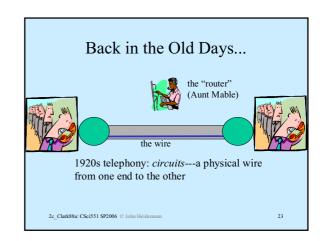
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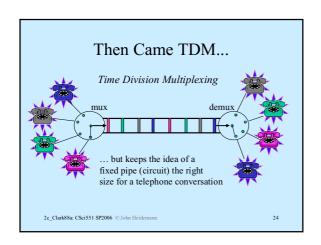
Main Goals

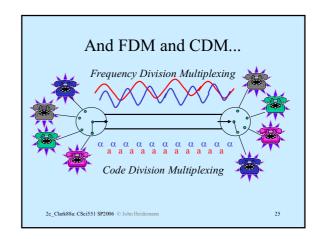
- · heterogeneous link-layers
 - why? already out there (built on telephone network), and wanted to let things evolve
 - some examples: wireless 802.11, Ethernet, X.25, optical, satellites, token rings, .
- multiplexing
 - packet switching fundamentally different from circuit switching

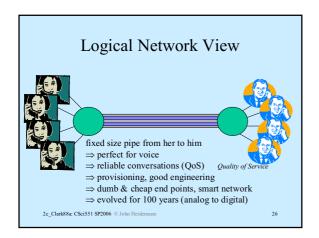
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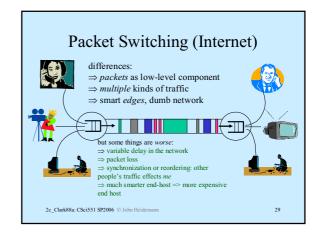
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Statistical Multiplexing Gain Assumptions: 1 Mb/s link user: 0.1Mb/s when transmitting, but 10% duty cycle • Circuit switching: can support 10 users, 100% reliable • Packet switching: with 35 users, probability that ≥10 are transmitting at the same time p = .1, N = 35, T = 10 $p = 1 - \sum_{i=1}^{t=T} {N \choose i} p_i^{i} (1-p_i)^{(N-i)}$

a. Robust to Failures • app should not see transient failures • what kinds of failures? - (talked about earlier today) • how does anything still work? - fate-sharing • soft-state • if we put state just at the ends, and we lose and end, we don't care - end-to-end argument • saving state at the end hosts - (replication)

b. Multiple Types of Service

- originally just NCP, but split to {TCP,UDP}/IP
- why?
 - to get different types of service
 - because different apps need different things from the
 - ex: with voice, you don't want relibability because you can just ask the other person to repeat

 downside of reliability:
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- · what?
- · made different protocols
- - reliable, end-to-end, connection oriented, byte
- UDP
 - unreliable, end-to-end, connectionless, packet

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Other protocols (than TCP/UDP/IP)?

- · RTP: Real Time Protocol
 - checksum, connectionless (?), standard frame of timing information
- RTSP: Real Time Streaming Protocol
- SCTP: Stream Control Transmission Protocol
 - streaming like TCP, but without reliability
- application protocols (above transport)
 - SIP: not really transport
 - HTTP, ...

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Non-TCP/UDP protocols

- RDP: Reliable Delivery Protocol message-based
 - allows out-of-order delivery
- RFC-908
- SCTP: Stream Control Transmission Protocol
 - intended for telephony signaling over IP
 - multiplexes multiple "streams" per connection
 - in-sequence per stream, out-of-sequence between streams
 - reliable

- · DCCP: Datagram Congestion Control Protocol
- TFRC: TCP-Friendly Rate
 - RFC-3448
 - we study TCP-friendliness later
- · XCP: Explicit Control
- Protocol

 - presented at SIGCOMM
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- in progress
- - defines how to do rate-control as a function of loss rate

- add congestion control to UDP
- Control

- high-speed streaming with explicit router rate feedback
- we'll study it later

- · classes:
 - interactive multimedia

 - interactive multimedia
 voice consensations
 [real time]
 gaming
 one-way multimedia
 video streaming
 bulk data transfer
 file transfer
 peer-to-peer
 e-mail
 interactive non-multimedia
 remote login
 - remote login
 chat
 - [e-commerce]

· requirements:

b.2. multiple *applications*

- interactive vs. background

high bandwidth

reliablity

• this is not a perfect, orthogonal

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c. multiple kinds of networks

- IP over X
- compare to integrated stacks:
 - ISO
 - ATM - cell phones (CPDP & WAP, maybe??)
 - fibre channel, Apple Desktop Bus, USB, Firewire
- but a counter example: SCSI and now SCSI over IP
- requirements of X:
 - reasonable size packets
 - but fragmentation and reassembly
 - reasonable reliablity
- but workarounds
- addressing • non-requirements of X:
 - reliable, in-order, broadcast, QoS, etc.

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d. other goals

- distributed management
 - some work, and today policy routing exists
 - but limitations (ex. address space portability)
- cost effective
 - today quite cheap
 - but for small devices? for keyboard?

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d. other other goals

- · effort to deploy end host
 - for him in '88: cost of implementing stack today: cost of administering machine
 - - much lower today (DHCP, etc.)
 but still lots of manual configuration "futzing"
- · accountability
 - tracking resource usage, money, identity of user and reprucussions of bad usage (spam, p2p file sharing, zombies and denial-of-service)
 - basically nothing then
 - today: not much, but today there is authentication (and identity) at many network connections, but not everywhere
 - · possibly a major focus of internet research will be look security

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Architecture and Implementation

- realization: an instance of the Internet class
 - him: 1200b/s modem vs. 1Mb/s LAN
 - today: from sensor net nodes with 8-bit CPUs and 20kb/s radios...to supercomputers
 - today: the Internet can't do X because it is Y
 - ex. can't do System Area Networks over IP because it's too slow, so we need Fibre Channel
 - alternative: build a fast Internet realization
 - corollary: not every realization is appropriate for every app
 - also: custom stack will get last 5% of performance, but is it worth it?

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TCP Alternative Choices

- byte stream vs. message stream
- · flow control
- · congestion control came later
- PSH flag
 - a weak record boundry
 - but the reason the Plan-9 people didn't use TCP

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Other Components of IP Success

- a good, free implementation
 - BSD Unix in the mid-80's
 - compare to OSI where impls were late
- a good API
 - BSD socket API
 - not perfect, but good
 - compare to OSes where Unix and Windows have very different APIs to open/rename/etc.

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Where are we now?

- this paper: The Internet in 1988 (!)
- much has changed since then (as discussed next in [Deering98a]
- what are the big challenges today?
 - think about this question... we'll come back to it

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Other questions/observations?

• XXX

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